New Directions in Numerical Linear Algebra and High Performance Computing: Celebrating the 70th Birthday of Jack Dongarra, Manchester, July 7-8, 2021 (virtual)

# Many eigenpair computation via Hotelling deflation 

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## Introduction

1. Problem statement ("many eigenpair computation"):

Let $\left(\lambda_{i}, v_{i}\right)$ be the eigenpairs of a $n \times n$ symmetric matrix $A$ with the eigenvalue ordering $\lambda_{1} \leq \lambda_{2} \leq \cdots \leq \lambda_{n}$. Compute the partial decomposition:

$$
A V_{n_{e}}=V_{n_{e}} \Lambda_{n_{e}}
$$

where $\Lambda_{n_{e}}=\operatorname{diag}\left(\lambda_{1}, \ldots, \lambda_{n_{e}}\right)$,
Interested in the cases where $n$ is "huge" and $n_{e}$ is "large"
2. Emerging applications:

- electronic structure calculations of Lithium-ion electrolyte,
- graphene,
- dyanmics analysis of viral capsids of supramolecular systems such as Zika and West Nile viruses (structural biology)


## Introduction

3. Existing approaches - an incomplete list:

- Full eigenvalue decomposition:
- LAPACK, ScaLAPACK, PLASMA, MAGMA
- ELPA (Eigenvalue Solves for Petaflop Applications), https://elpa.mpcdf.mpg.de/
- EigenExa, https://www.r-ccs.riken.jp/labs/lpnctrt/projects/eigenexa/
- ...
- QDWHeig [Sukkari, Ltaief, Keyes at KAUST]
- SLATE [Gates et al, U. of Tennessee]

Stable, but expensive, $O\left(n^{2}\right)$ storage and $O\left(n^{3}\right)$ flops

- "Spectrum slicing:"
- SLEPc, https://slepc.upv.es/
- EVSL, http://www.cs.umn.edu/~saad/software
- FEAST, http://www.ecs.umass.edu/~polizzi/feast/
- z-Pares, http://zpares.cs.tsukuba.ac.jp/
- ...
- SISLICE [Williams-Young and Yang, LBNL]

Scalable, but issues with duplicate/missing eigenvalues between slices, ...

## Introduction

4. Using Lanczos or any other subspace projection methods for many eigenpairs are challenging - numerically and computationally:

- needs large subspace (memory) $M$, e.g., $m=2 n_{e}$
- require (internal) locking (deflation) to avoid danger of converging again to the same eigenvalues ( $O\left(\mathrm{~nm}^{2}\right.$ ) flops)


## Introduction

5. Many eigenpair computation is challenging even when vast computational resources are available.
Case demo: EVSL (http://www.cs.umn.edu/~saad/software)






Left: Dengue virus model (www.rcsb.org/structure/4cct). $n=307,260$ Middle: DOS sliced 5, 076 eigenvalues into 20 subintervals (top) and the relative residual norms of eigenpairs (bottom). $\left\|\widehat{V}^{T} \widehat{V}-I\right\|_{F}=O\left(10^{-6}\right)$. Right: Degrees of filter polynomials (top) CPU timing (bot.) for each slices

## Introduction

6. This talk is about an on-going project on

Lanczos + Hotelling deflation + Communication-avoiding for computing many eigenpairs.

## Rest of the talk

I. Hotelling deflation = Explicit External Deflation (EED)
II. Backward stability of EED
III. Communication-avoiding algorithm for MPK (CA-MPK)
IV. Eigensolver $s$ TRLED ( $=$ TRLan + EED + CA-MPK)
V. Concluding remarks

## Joint work with

- Jack J. Dongarra, Univ of Tennessee
- Chao-Ping Lin, UC Davis
- Ding Lu, Univ of Kentucky
- Ichitaro Yamazaki, Sandia National Labs.


## I. Explicit External Deflation (EED)

1. Hotelling deflation (EED $=$ Explicit External Deflation)

Let $A V=V \Lambda$ be the eigen-decomposition of $A$, partition

$$
V=\left[\begin{array}{ll}
V_{k} V_{n-k}
\end{array} \quad \text { and } \quad \Lambda=\left[\begin{array}{ll}
\Lambda_{k} & \\
& \Lambda_{n-k}
\end{array}\right]\right.
$$

and define

$$
\widehat{A}=A+V_{k} \Sigma_{k} V_{k}^{T}
$$

where $\Sigma_{k}=\operatorname{diag}\left(\sigma_{1}, \sigma_{2}, \ldots, \sigma_{k}\right)$ are shifts. Then
(a) The eigenvalues of $\widehat{A}$ are

$$
\lambda_{i}(\widehat{A})= \begin{cases}\lambda_{i}+\sigma_{i} & \text { for } 1 \leq i \leq k \\ \lambda_{i} & \text { for } k+1 \leq i \leq n\end{cases}
$$

(b) $A$ and $\widehat{A}$ have the same eigenvectors

Therefore, one can use proper shifts $\sigma_{i}$ to move "computed" eigenvalues away, and then compute the next batch of "favorite" eigenvalues for an eigensolver.

## I. Explicit External Deflation

2. Governing equations of the exact EED process: for $j=1,2, \ldots$,

$$
\begin{aligned}
A_{j} & =A_{j-1}+\sigma_{j} v_{j} v_{j}^{T}=A+V_{j} \Sigma_{j} V_{j}^{T} \quad \text { and } A_{0}=A . \\
A_{j} v_{j+1} & =\lambda_{j+1} v_{j+1} \\
A_{j} V_{j} & =V_{j}\left(\Lambda_{j}+\Sigma_{j}\right)
\end{aligned}
$$

with initial $A v_{1}=\lambda_{1} v_{1}$, where $\Sigma_{j}=\operatorname{diag}\left(\sigma_{1}, \ldots, \sigma_{j}\right)$.

## I. Explicit External Deflation

## 3. Benefits of EED:

- Easily incooperated into existing eigensolvers, such as TRLan and ARPACK.
- Small projection subspace dimensions (core memory requirements) even for many eigenpairs.
- No need to explicitly (re)-orthogonalize the projection subspace to the computed eigenvectors.
- Accelerated convergence with warm start for $A_{j}$.
- Straightforward extension to generalized symmetric eigenproblem $A v=\lambda B v$ :

$$
\left(A+\sigma B V_{k} V_{k}^{T} B\right) x=\lambda B x
$$

- Readily exploit structures of $A$ and $B$, for example, see EED for the linear response eigenvalue problem [Bai-Li-Lin'17].
- Extensible to other eigenvalue-type problems, such as SVD, sparse PCA.
- Nonlinear eigenvector problem (NEPv) analogy: level shifting


## I. Explicit External Deflation

4. Two key issues of EED:
(a) Numerical linear algebra issue:

Numerical stability with approximate eigenvectors $\widehat{V}_{k}$ :

$$
\widehat{A}_{j}=A+\widehat{V}_{j} \Sigma_{j} \widehat{V}_{j}^{T}
$$

(b) High performance computing issue:

Cost of matrix powers kernel (MPK) for generating Krylov subspace:

$$
\left[p_{0}\left(\widehat{A}_{j}\right) v_{0}, p_{1}\left(\widehat{A}_{j}\right) v_{0}, \ldots, p_{s}\left(\widehat{A}_{j}\right) v_{0}\right]
$$

where $\left\{p_{k}(\cdot)\right\}$ are recursively defined polynomials.

## II. Backward stability of EED

1. Mixed messages from previous work:
[Wilkinson'65], [Parlett'82], [Saad'89], [Jang/Lee'06], ...
2. Governing equations of inexact EED:

$$
\begin{aligned}
\widehat{A}_{j} & =\widehat{A}_{j-1}+\sigma_{j} \widehat{v}_{j} \widehat{v}_{j}^{T}=A+\widehat{V}_{j} \Sigma_{j} \widehat{V}_{j}^{T} \\
\widehat{A}_{j} \widehat{v}_{j+1} & =\widehat{\lambda}_{j+1} \widehat{v}_{j+1}+\eta_{j+1} \\
\widehat{A}_{j} \widehat{V}_{j} & =\widehat{V}_{j}\left(\widehat{\Lambda}_{j}+\Sigma_{j}\right)+\widehat{V}_{j} \Sigma_{j} \Phi_{j}+E_{j}
\end{aligned}
$$

where $\widehat{A}_{0}=A$,

$$
\begin{aligned}
\left\|\eta_{j+1}\right\| & \leq \operatorname{tol} \cdot\|A\|, \\
\Phi_{j} & =\operatorname{utri}\left(\widehat{V}_{j}^{T} \widehat{V}_{j}-I_{j}\right), \\
E_{j} & =\left[\eta_{1}, \eta_{2}, \ldots, \eta_{j}\right] .
\end{aligned}
$$

and $t o l$ is prescribed relative residual tolerance.

## II. Backward stability of EED

3. Metrics of backward stability for computed eigenpairs $\left(\widehat{\Lambda}_{j+1}, \widehat{V}_{j+1}\right)$ with relative residual tolerance tol:

- The loss of orthogonality

$$
\omega_{j+1}=\left\|\widehat{V}_{j+1}^{T} \widehat{V}_{j+1}-I\right\|_{F}=O(t o l)
$$

- The symmetric backward error norm

$$
\delta_{j+1}=\min _{\Delta \in \mathcal{H}}\|\Delta\|_{F}=O(t o l \cdot\|A\|),
$$

where

$$
\mathcal{H}=\left\{\Delta \mid(A+\Delta) Q_{j+1}=Q_{j+1} \widehat{\Lambda}_{j+1}, \Delta=\Delta^{T}, Q_{j+1}=\operatorname{orth}\left(\widehat{V}_{j+1}\right)\right\}
$$

## II. Backward stability of EED

4. Two key quantities associated with the shifts

- Spectral gap

$$
\gamma_{j} \equiv \min _{\lambda \in \mathcal{J}_{j+1}, \theta \in \mathcal{J}_{j}}|\lambda-\theta|,
$$

where $\mathcal{J}_{j+1}=\left\{\widehat{\lambda}_{1}, \ldots, \widehat{\lambda}_{j}, \widehat{\lambda}_{j+1}\right\}$ is the set of computed eigenvalues, and $\mathcal{J}_{j}=\left\{\widehat{\lambda}_{1}+\sigma_{1}, \ldots, \widehat{\lambda}_{j}+\sigma_{j}\right\}$ is the set of computed eigenvalues with shifts.

- Shift-gap ratio

$$
\tau_{j} \equiv \frac{1}{\gamma_{j}} \cdot \max _{1 \leq i \leq j}\left|\sigma_{i}\right|
$$

5. Theorem.

Under mild assumptions, if

$$
\begin{equation*}
\gamma_{j}^{-1}\|A\|=O(1) \quad \text { and } \quad \tau_{j}=O(1) \tag{1}
\end{equation*}
$$

then

$$
\omega_{j+1}=O(t o l) \quad \text { and } \quad \delta_{j+1}=O(t o l \cdot\|A\|)
$$

6. Rule of Thumb:
dynamical choice of shifts $\sigma_{j}$ to satisfy the conditions (1).

## II. Backward stability of EED

7. Numerics of TRLED $=$ TRLan + EED:

- Test matrices:

| matrix | $n$ | $\left[\lambda_{\min }, \lambda_{\max }\right]$ | $\left[\lambda_{\text {low }}, \lambda_{\text {upper }}\right]$ | $n_{e}$ |
| :---: | :---: | :---: | :---: | :---: |
| Laplacian | 40,000 | $[0,7.9995]$ | $[0,0.07]$ | 205 |
| worms20 | 20,055 | $[0,6.0450]$ | $[0,0.05]$ | 289 |
| SiO | 33,401 | $[-1.6745,84.3139]$ | $[-1.7,2.0]$ | 182 |
| Si34H36 | 97,569 | $[-1.1586,42.9396]$ | $[-1.2,0.4]$ | 310 |
| Ge87H76 | 112,985 | $[-1.214,32.764]$ | $[-1.3,-0.0053]$ | 318 |
| Ge99H100 | 112,985 | $[-1.226,32.703]$ | $[-1.3,-0.0096]$ | 372 |

## II. Backward stability of EED

7. Numerics of TRLED $=$ TRLan + EED, cont'd

- Results:

| matrix | $\widehat{n}_{e}$ | $j_{\max }$ | $\omega_{\widehat{n}_{e}}$ | $\left\\|R_{\widehat{n}_{e}}\right\\|_{\mathrm{F}} /$ Anorm | CPU time (sec.) |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | TRLED | TRLan |  |  |
| Laplacian | 205 | 60 | $1.93 \cdot 10^{-8}$ | $6.33 \cdot 10^{-8}$ | 66.5 | 86.0 |
| worms20 | 289 | 86 | $2.63 \cdot 10^{-8}$ | $7.24 \cdot 10^{-8}$ | 57.3 | 74.8 |
| Si0 | 182 | 41 | $2.33 \cdot 10^{-8}$ | $4.71 \cdot 10^{-8}$ | 42.4 | 47.1 |
| Si34H36 | 310 | 72 | $3.41 \cdot 10^{-8}$ | $7.50 \cdot 10^{-8}$ | 309.9 | 310.4 |
| Ge87H76 | 318 | 66 | $4.08 \cdot 10^{-8}$ | $8.50 \cdot 10^{-8}$ | 388.7 | 421.0 |
| Ge99H100 | 372 | 74 | $3.65 \cdot 10^{-8}$ | $7.63 \cdot 10^{-8}$ | 501.1 | 533.4 |

- EED profile of Ge99H100:




## II. Backward stability of EED

7. Numerics of TRLED $=$ TRLan + EED, cont'd

- Observations:
(1) All desired eigenvalues are successfully computed: $n_{e}=\widehat{n}_{e}$
(2) all computed eigenpairs are backward stable:

$$
\omega_{n_{e}}=O(t o l) \quad \text { and } \quad \delta_{n_{e}} \approx\left\|R_{n_{e}}\right\|=O(t o l \cdot\|A\|)
$$

(3) EED does not slow down, in fact, slightly faster, partially due to smaller "internal" memory usage and warm start.
8. With proper choice of shifts, Hotelling deflation (EED) would not compromise numerical stability of an eigensolver.
[Lin, Lu and Bai, arXiv:2105.01298, May 2021]

## III. Communication-avoiding algorithm for computing MPK

1. Sparse-plus-low-rank matrix powers kernel (MPK)

$$
\left[p_{0}(B) v, p_{1}(B) v, \ldots, p_{s}(B) v\right]
$$

with

$$
B=A+\sigma U_{k} U_{k}^{T}=\text { sparse }+ \text { low rank }
$$

and $A U_{k}=U_{k} \Lambda_{k}, U_{k}^{T} U_{k}=I_{k}$ and $p_{j}(\cdot)$ are polynomials defined recursively.
2. For simplicity, consider polynomials $\left\{p_{j}(\cdot)\right\}$ in the monomial basis and compute the MPK:

$$
\begin{aligned}
{\left[p_{0}(B) v, p_{1}(B) v, \ldots, p_{s}(B) v\right] } & =\left[x, B x, B^{2} x, \ldots, B^{s} x\right] \\
& \equiv\left[x^{(0)}, x^{(1)}, x^{(2)}, \ldots, x^{(s)}\right]
\end{aligned}
$$

3. Standard MPK algorithm for computing $x^{(1)}, x^{(2)}, \ldots$

1: $x^{(0)}=x$;
2: for $j=1: s$ do
3: $\quad x^{(j)}=B x^{(j-1)}=A x^{(j-1)}+\sigma\left(U_{k}\left(\underline{U_{k}^{T} x^{(j-1)}}\right)\right)$
4: end for

## III. Communication-avoiding algorithm for computing MPK

4. Performance example

- $B=A+\sigma U_{k} U_{k}^{T}$, where $A$ is 2D Laplacian and $U_{k}$ are eigenvectors
- MPK $V_{s}=\left[p_{0}(B) v_{0}, p_{1}(B) v_{0}, \ldots, p_{s}(B) v_{0}\right]$
- Timing of the standard algorithm in MATLAB on a desktop - blue curve.

- Q: what's the red curve?


## III. Communication-avoiding algorithm for computing MPK

5. An algorithm has two costs:
```
arithmetic (flops) + movement of data (communication)
```

6. Communication is the bottleneck on modern architectures.
7. Q: How to exploit the sparse-plus-low-rank matrix structure to reduce communication cost?

Ans.: use a specialized communication-avoiding algorithm developed by

- Leiserson-Rao-Toledo'97 ("out-of-core") and
- Knight-Carson-Demmel'13 ("exploiting data sparsity")


## III. Communication-avoiding algorithm for computing MPK

8. Communication-Avoiding (CA) algorithm for the MPK

1: $x^{(0)}=x$
2: $b_{0}=U_{k}^{T} x^{(0)}$
3: $W_{k}^{(j)}=\Lambda_{k}^{j}+\sigma \sum_{i=1}^{j} \Lambda_{k}^{i-1}\left(\Lambda_{k}+\sigma\right)^{j-i}$ for $j=1: s-1$,
4: $b_{j}=W_{k}^{(j)} b_{0}$ for $j=1: s-1$
5: $\left[c_{0}, c_{1}, \ldots, c_{s-1}\right]=U_{k}\left[b_{0}, b_{1}, \ldots, b_{s-1}\right]$
6: for $j=1: s$ do
7: $\quad x^{(j)}=A x^{(j-1)}+\sigma c_{j-1}$
8: end for
9. Benefits of CA-MPK algorithm:

- Reduced flops

$$
\begin{aligned}
\text { flops }_{\text {std }} & =n n z A \cdot s+n k s+n k s \\
\text { flops }_{\mathrm{ca}} & =n n z A \cdot s+n k s+n k+O\left(k^{2} s\right)
\end{aligned}
$$

- Reduced movement of data: $U_{k}$ is only accessed twice.


## III. Communication-avoiding algorithm for computing MPK

10. Performance example, cont'd:

- $B=A+\sigma U_{k} U_{k}^{T}$, where $A$ is 2D Laplacian and $U_{k}$ are eigenvectors
- MPK $V_{s}=\left[p_{0}(B) v_{0}, p_{1}(B) v_{0}, \ldots, p_{s}(B) v_{0}\right]$
- Timing in MATLAB
- Standard algorithm - blue curve
- CA algorithm - red curve



## III. Communication-avoiding algorithm for computing MPK

10. Performance example, cont'd:

- $B=A+\sigma U_{k} U_{k}^{T}$, where $A$ is 2D Laplacian and $U_{k}$ are eigenvectors
$-\mathrm{MPK} V_{s}=\left[p_{0}(B) v_{0}, p_{1}(B) v_{0}, \ldots, p_{s}(B) v_{0}\right]$
- Timing in one node (32 cores) of Cori (NERSC)



11. Rounding error analysis of CA-MPK

## IV. Lanczos algorithm with EED and MPK

1. Lanczos algorithm

- Lanczos process

$$
A Q_{m}=Q_{m} T_{m}+\beta_{m} q_{m+1} e_{m}^{T}
$$

where

$$
\operatorname{span}\left\{Q_{j}\right\}=\operatorname{span}\left\{q, A q, \ldots, A^{m-1} q\right\} \quad \text { (Krylov subspace) }
$$

- Rayleigh-Ritz approximation

$$
\begin{aligned}
T_{m} x_{i} & =\theta_{i} x_{i} \\
\left(\lambda_{i}, v_{i}\right) & \approx\left(\theta_{i}, Q_{m} x_{i}\right)
\end{aligned}
$$

2. Lacnzos algorithm is efficient for computing a few exterior eigenvalues (and eigenvectors).
3. Two main kernels

- Matrix-Vector multiply (SpMV) for generating Krylov subspace: $A q$
- Re-orthogonalization for maintaining orthonormal basis vectors: $Q_{j}$

4. Two variants of Lanczos method:

- Thick-restart Lanczos (TRLan) - control $m$ size
- s-step Lanczos (s-Lanczos) - reduce communication cost by using MPK


## IV. Lanczos algorithm with EED and MPK

5. sTRLED $=s$-step-TRLan + EED + CA-MPK
6. Preliminary results on strong-parallel scaling of sTRLED on multi-processor for Si87H76:
$n=240,369, n_{e}=700$
computing 100 eigenvalues at a time with $m=200$ and $s=5$


[Bai, Dongarra, Lu, Yamazaki, IPDPS19]

## V. Concluding remarks

1. Many eigenpairs computation in emerging applications, a challenging problem even when vast computational resources are available.
2. Two techniques discussed in this talk:

- Explicit external deflation (EED) for reliably moving away computed eigenpairs,
- a communication-avoiding matrix powers kernel (CA-MPK) for fast sparse-plus-low-rank MPK

3. The capability of being able to efficiently compute large number of eigenvalues will not just be appealing, but also mandatory for the next generation of eigensolvers.
